MEM6804 Modeling and Simulation for Logistics & Supply Chain 物流与供应链建模与仿真

Theory

Lecture 4: Random Variate Generation

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2 Random Number Generation

- Definition
- Pseudo-Random Numbers
- Linear Congruential Generator
- ► More Sophisticated RNGs
- ► Tests for Random Numbers

3 Random Variate Generation

- Inverse-Transform Technique
- Acceptance-Rejection Technique
- Other Ad-Hoc Methods
- Generating Poisson Process



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- When simulating a system, we often need to generate random variates (e.g., interarrival time, service time) from all kinds of distributions (e.g., exponential distribution, arbitrary empirical distribution).



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 - To better understand the randomness in stochastic simulation.
 - Be alert to some inadequate random variate generator.
- To produce a sequence of random variates from a given distribution (of a random variable):
 - Start with random variates from Unif(0, 1) (called random numbers).
 - 2 All random variates with given distribution are "transformed" from random numbers.

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 - If $U \sim \text{Unif}(0, 1)$, then $\mathbb{E}[U] = \frac{1}{2}$, $\text{Var}(U) = \frac{1}{12}$, and its pdf is $f(u) = \begin{cases} 1, & 0 \le u \le 1, \\ 0, & \text{otherwise.} \end{cases}$

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- Statistical Properties
 - Uniformity: Each value on [0, 1] has equal likelihood.
 - Independence: Implies no correlation between successive numbers.



Definition

Uniformity

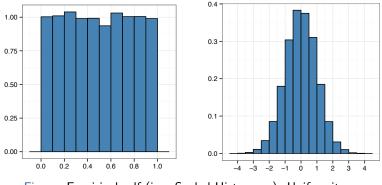
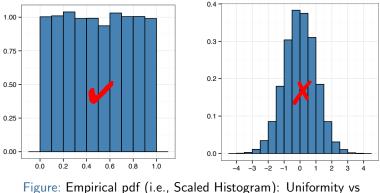


Figure: Empirical pdf (i.e., Scaled Histogram): Uniformity vs Nonuniformity (from ZHANG Xiaowei)



• Uniformity



Nonuniformity (from ZHANG Xiaowei)



Independence

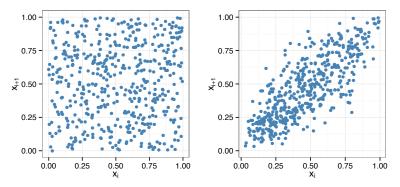


Figure: Scatter Plot: Uncorrelated vs Correlated (from ZHANG Xiaowei)



► Definition

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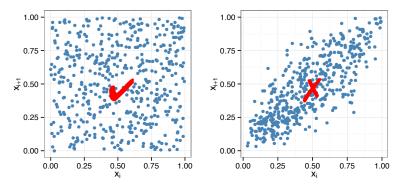


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- "Pseudo" means false
 - Generating random numbers by a known method removes true randomness.
 - The set of pseudo-random numbers can be repeated.
- Goal: To produce a sequence of numbers in [0, 1] that imitates the ideal properties of random numbers.
 - Statistical properties are the most important.
 - True randomness is not the first priority.



- Properties of a good random number generator (RNG):
 - Pass statistical tests.
 - 2 Solid theoretical support.
 - I Fast.
 - Sufficiently long cycle (period).
 - **5** Portable to different computers.
 - Replicable.



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- Techniques for RNG:
 - Linear Congruential Generator (LCG)
 - Combined LCG
 - Multiple Recursive Generator (MRG)



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• The selection of the values for a, c, m, and x_0 drastically affects the statistical properties and the cycle length $M \ge \# \xi \not\in \mathcal{F} \not\cong \mathcal{F}$

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• Try https://xiaoweiz.shinyapps.io/randNumGen for different parameters.

- An actual use of LCG (Lewis et al. 1969): $a = 7^5$, c = 0, $m = 2^{31} 1 = 2,147,483,647$ (a prime number).
 - It adopts $u_i = \frac{x_i}{m+1}$.
 - It passes many of the standard statistical tests.
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- Note: By letting modulus *m* be a power of 2 (or close), the modulo operation can be conducted efficiently, since most digital computers use a binary representation of numbers.
- As computing power has increased, LCG is not adequate nowadays; more sophisticated RNGs are used in practice.



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3 Let $x_{j+1} = (x_{1,j+1} - x_{2,j+1}) \mod (m_1 - 1)$. (*Remark*: mod uses floored division, i.e., $y \mod m = y - m\lfloor \frac{y}{m} \rfloor$.)

4 Return

$$u_{j+1} = \begin{cases} \frac{x_{j+1}}{m_1}, & \text{if } x_{j+1} > 0, \\ \frac{m_1 - 1}{m_1}, & \text{if } x_{j+1} = 0. \end{cases}$$

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It has cycle length $(m_1-1)(m_2-1)/2 \approx 2 \times 10^{18}$. If $M_1 = 10^{18}$

• Multiple Recursive Generator (MRG): Extends LCG by using a higher-order recursion:

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 - It has cycle length $\approx 3\times 10^{57}$, which is enormous.
 - If you could generate one billion (10⁹) pseudo-random numbers per second, then it would take longer than the age of the universe to exhaust the period of MRG32k3a!



[†]MRG32k3a or its adaptation is one of the RNGs used in MATLAB, R, SAS, Arena, etc.

- Tests based on generated sequences of numbers.
 - Frequency Test for uniformity (discussed in next lecture)
 - Kolmogorov-Smirnov test (柯尔莫哥洛夫-斯米尔诺夫检验)
 - chi-square test (χ^2 test, 卡方检验)
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- Fortunately, the well-known RNGs which are widely used in simulation softwares and languages have been extensively tested and validated.
- Be careful when the RNG at hand is not explicitly known or documented!
 - Even RNGs that have been used for years in popular commercial softwares (e.g., Excel, Visual Basic), have been found to be inadequate (L'Ecuyer 2001).

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- Goal: Produce random variates from a given probability distribution (e.g. exponential, Poisson, etc.).
- Widely-used techniques[†]
 - Inverse-transform technique (generic)
 - Acceptance-rejection technique (generic)
 - Other ad-hoc methods for some specific distributions

^T Methods introduced in this lecture are exact; there are also approximation methods such as MCMC.^{HM IMO TOWE UNIV}



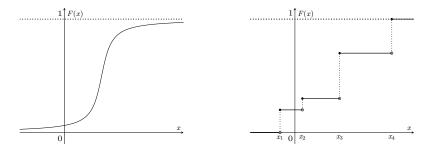


Figure: Continuous Random Variable

Figure: Discrete Random Variable



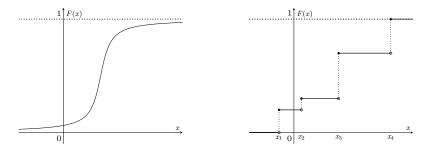


Figure: Continuous Random Variable

Procedures

Figure: Discrete Random Variable



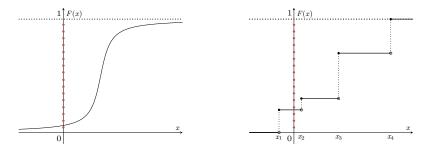


Figure: Continuous Random Variable

Figure: Discrete Random Variable

• Procedures

1 Generate (as needed) random numbers (on vertical axis).



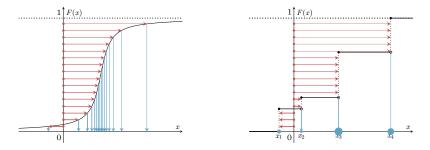


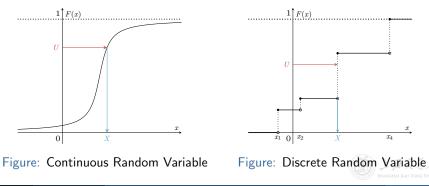
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- Procedures
 - **1** Generate (as needed) random numbers (on vertical axis).
 - 2 Map inversely to points on horizontal axis, which are the desired random variates from F(x).

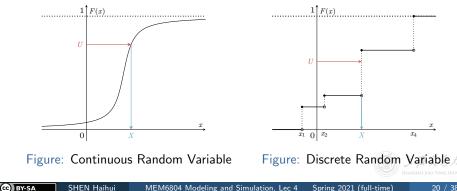
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- If $U \sim \text{Unif}(0, 1)$, then $F^{-1}(U)$ has the same distribution as X. i.e.,

$$\mathbb{P}(F^{-1}(U) \le x) = \mathbb{P}(U \le F(x)) = F(x).$$



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- It can be used to sample from all (in principle) discrete distributions, e.g.,
 - discrete uniform
 - geometric

SHEN Haihui

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arbitrary empirical distribution



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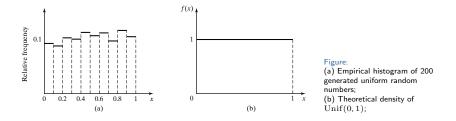
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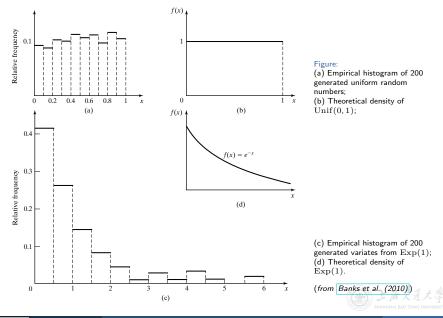
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- Remark: $1 U \sim \text{Unif}(0, 1) \Longrightarrow -\frac{1}{\lambda} \ln(U)$ is sufficient.
- Numerical test for Exp(1) in Excel.

Generate 200 random numbers.

Obtain 200 random variates via the inverse function.







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Discrete Distribution

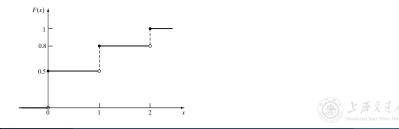
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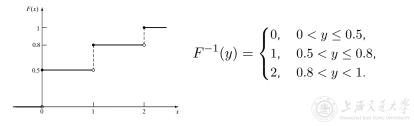
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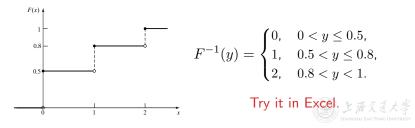
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- Acceptance-rejection technique is also useful for generating a *non-stationary Poisson process* (more details later).



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 - Whereas there exists a one-to-one mapping for the inverse-transform method.



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 - U itself does not have the desired distribution.
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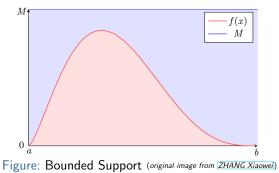
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• For
$$1/4 \le x \le 1$$
,
 $\mathbb{P}\{U \le x | U \ge 1/4\} = \frac{\mathbb{P}\{U \le x \text{ and } U \ge 1/4\}}{\mathbb{P}\{U \ge 1/4\}} = \frac{x - 1/4}{3/4}$,

which is exactly the desired CDF of $X \sim \text{Unif}(1/4, 1)$.



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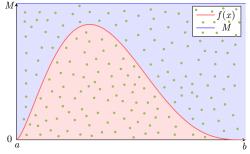


Figure: Bounded Support (original image from ZHANG Xiaoweil)

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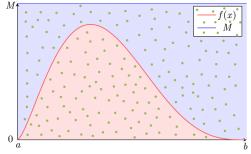


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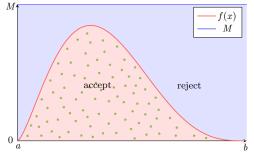


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• The acceptance rate is $\mathbb{P}\{Z < f(Y)\} = \frac{1}{(b-a)M}$.

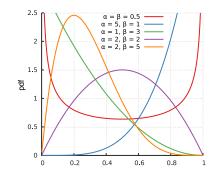


Bounded Support

• Goal: Generate random variates from $Beta(\alpha, \beta)$, where the density is $f(x) = \frac{x^{\alpha-1}(1-x)^{\beta-1}}{B(\alpha, \beta)}$, $x \in [0, 1]$.



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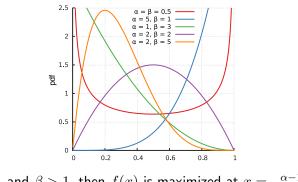




Beta from Uniform

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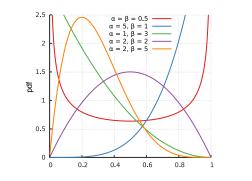
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• If $\alpha > 1$ and $\beta > 1$, then f(x) is maximized at $x = \frac{\alpha - 1}{\alpha + \beta - 2}$ and the maximum is $M = \frac{(\alpha - 1)^{\alpha - 1}(\beta - 1)^{\beta - 1}}{(\alpha + \beta - 2)^{\alpha + \beta - 2}B(\alpha, \beta)}$.

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• The acceptance rate is $\frac{1}{(b-a)M} = \frac{1}{(1-0)M} = \frac{1}{M}$.

SHEN Haihui

Beta from Uniform

• Generate random variates from X, whose density f(x) is upper bounded by Mg(x), where g(x) is *instrumental* density.

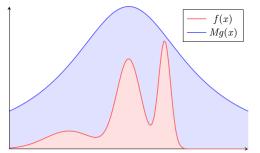


Figure: Unbounded Support (original image from ZHANG Xiaowei)



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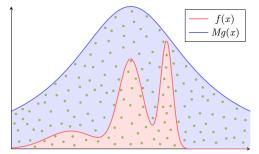


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• Generate random variate pairs (y_1, z_1) , (y_2, z_2) , ..., from Uniform $\{(y, z) : y \in \text{support of } g(\cdot), 0 \le z \le Mg(y)\}$.



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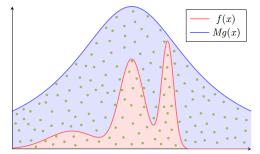


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 - y_i from $Y \sim g(\cdot)$, z_i from $Z \sim \text{Unif}(0, Mg(y_i))$ (why?)



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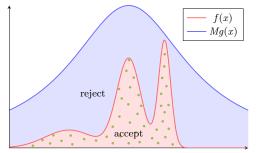


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- 2 Accept the pair if $z_i < f(y_i)$ and output y_i as random variate from X with density f(x).

- Y conditioned on the event $\{Z < f(Y)\}$ has the same distribution as X, i.e., having density f(x).
 - Let Θ denote $\{(y, z) : y \in \text{support of } g(\cdot), 0 \le z \le Mg(y)\}.$
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 $\mathbb{P}\{Y \leq x | Z < f(Y)\}$



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• The acceptance rate is $\mathbb{P}\{Z < f(Y)\} = \frac{1}{\Theta \text{ area}} = \frac{1}{\int_{-\infty}^{\infty} Mg(y) \mathrm{d}y} = \frac{1}{M \int_{-\infty}^{\infty} g(y) \mathrm{d}y} = \frac{1}{M} \cdot (1 + f(y))$

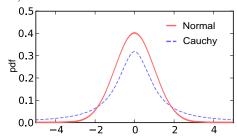
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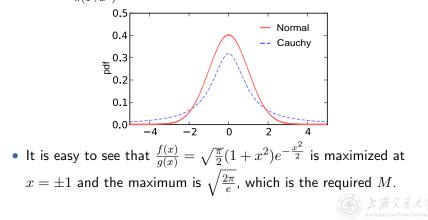


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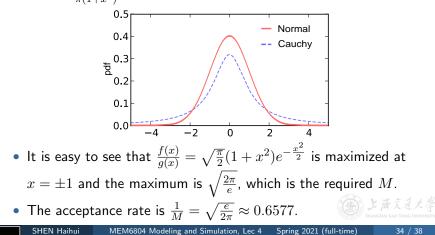
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Normal from Cauchy

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- Box–Muller method for $\mathcal{N}(0, 1)$ random variates:
 - **1** Generate u_1 and u_2 independently from Unif(0, 1).
 - 2 Let $z_1 = \sqrt{-2 \ln u_1} \cos(2\pi u_2)$ and $z_2 = \sqrt{-2 \ln u_1} \sin(2\pi u_2)$.



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- Intuition:
 - For two independent $\mathcal{N}(0,1)$ RVs Z_1 and Z_2 ,

$$Z_1^2, Z_2^2 \sim \chi_1^2, \ Z_1^2 + Z_2^2 \sim \chi_2^2.$$

- $X \sim \operatorname{Exp}(1/2) \iff X \sim \chi_2^2$.
- $-2\ln u_1$ is a random variate from $\operatorname{Exp}(1/2)$ (and thus χ_2^2).
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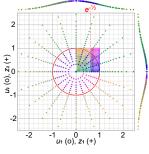
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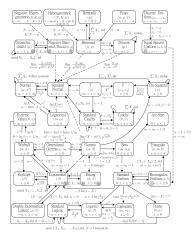


Figure: Relationships Among 35 Distributions (from Song (2005))

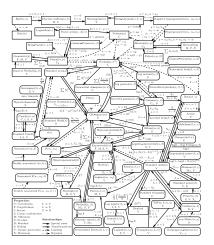


Figure: Relationships Among 76 Distributions (from [Leemis & McQueston (2008)])



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Nonhomogeneous Poisson process with rate (intensity) function λ(t):

$$N(t+h) - N(t) \sim \text{Poisson}(m(t+h) - m(t)),$$
 where $m(t) = \int_0^t \lambda(s) \mathrm{d}s.$



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 - 2 Generate x from Exp(λ*), and let t ← t + x (this is the arrival time of the stationary Poisson process with rate λ*).



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 - Generate a stationary Poisson arrival process at the fastest rate $\lambda^* = \max_t \lambda(t)$.
 - But "accept" only a portion of arrivals, thinning out just enough to get the desired time-varying rate.
- Algorithm:
 - **1** Set t = 0 and i = 1.
 - 2 Generate x from Exp(λ*), and let t ← t + x (this is the arrival time of the stationary Poisson process with rate λ*).
 - **3** Generate random number u (from Unif(0, 1)). If $u \leq \lambda(t)/\lambda^*$, then $s_i = t$ and $i \leftarrow i + 1$.

- To generate nonhomogeneous Poisson process with rate function $\lambda(t)$, one can use the acceptance-rejection method (which is also called *thinning* in this context).
- Idea behind thinning:
 - · Generate a stationary Poisson arrival process at the fastest rate $\lambda^* = \max_t \lambda(t).$
 - But "accept" only a portion of arrivals, thinning out just enough to get the desired time-varying rate.
- Algorithm:

 - **2** Generate x from $\text{Exp}(\lambda^*)$, and let $t \leftarrow t + x$ (this is the arrival time of the *stationary* Poisson process with rate λ^*).
 - **3** Generate random number u (from Unif(0, 1)). If $u \leq \lambda(t)/\lambda^*$, then $s_i = t$ and $i \leftarrow i+1$.
 - 4 Go to Step 2.

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